

# COMPUTER neuvories

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Position of IP and other network-layer protocols in TCP/IP protocol suite



- IPv4 is an unreliable datagram protocol—a best-effort delivery service.
- The term best-effort means that IPv4 packets can be corrupted, be lost, arrive out of order, or be delayed, and may create congestion for the network.
- If reliability is important, IPv4 must be paired with a reliable transport-layer protocol such as TCP

#### Example of Best Effort Delivery Service

- The post office does its best to deliver the regular mail but does not always succeed.
- If an unregistered letter is lost or damaged, it is up to the sender or would-be recipient to discover this.
- The post office itself does not keep track of every letter and cannot notify a sender of loss or damage of one





• IPv4 is also a connectionless protocol that uses the datagram approach.

- Packets used by the IP are called datagrams .
- A datagram is a variable-length packet consisting of two parts: header and payload (data).
- The header is 20 to 60 bytes in length and contains information essential to routing and delivery.

#### Datagrams

#### Packet at various Layers



#### 20-60 B



20– 20–60 bytes			5,535 bytes	→	VER: version number HLEN: header length	
	Header		Payload	byte: 8 bits		
	a. IP datagram			Flags D M		
_	0 4		8 <b>16</b>		31	
	VER 4 bits	HLEN 4 bits	Service type 8 bits		Total length 16 bits	
		Identification 16 bits		Flags 3 bits	Fragmentation offset 13 bits	
	Time-to-live 8 bits		Protocol 8 bits		Header checksum 16 bits	
	Source IP address (32 bits)					
Destination IP address (32 bits)					32 bits)	
	, Options + padding (0 to 40 bytes)				ng	

b. Header

#### **Version Number Field**

The 4-bit version number (VER) field defines the version of the IPv4 protocol, which, obviously, has the value of 4.





- The 4-bit header length (HLEN) field defines the total length of the datagram header in 4byte words.
- The receiver needs to multiply the value of this field by 4 to find the total length.





b. Header



- ✓ This 16-bit field defines the total length (header plus data) of the IP datagram in bytes.
- ✓ This field helps the receiving device to know when the packet has completely arrived.

#### Length of Data in IP Datagram

#### Length of data = total length – (HLEN) $\times 4$

# Why we need Total Length field

- However, there are occasions in which the datagram is not the only thing encapsulated in a frame; it may be that padding has been added.
- For example, the Ethernet protocol has a minimum and maximum restriction on the size of data that can be encapsulated in a frame (46 to 1500 bytes).
- If the size of an IPv4 datagram is less than 46 bytes, some padding will be added to meet this requirement.
- In this case, when a machine decapsulates the datagram, it needs to check the total length field to determine how much is really data and how much is padding.

### Identification, Flags, and Fragmentation Offset

Flags a. IP datagram D 8 **16** 31 0 **HLEN** Service type Total length VER 8 bits 16 bits 4 bits 4 hits Identification Fragmentation offset Flags 13 bits 16 bits 3 bits Header checksum lime-to-live Protocol 16 bits 8 bits 8 bits Source IP address (32 bits) Destination IP address (32 bits) Options + padding (0 to 40 bytes)

b. Header

# Fragmentation

- A datagram can travel through different networks. Each router decapsulates the IP datagram from the frame it receives, processes it, and then encapsulates it in another frame.
- For example, if a router connects a LAN to a WAN, it receives a frame in the LAN format and sends a frame in the WAN format.

### Fragmentation





# Maximum Transfer Unit (MTU)

- Each link-layer protocol has its own frame format.
- One of the features of each format is the maximum size of the payload that can be encapsulated.
- In other words, when a datagram is encapsulated in a frame, the total size of the datagram must be less than this maximum size,



MTU: Maximum size of frame payload

### Fragmentation

- The value of the MTU differs from one physical network protocol to another.
- For example, the value for a LAN is normally 1500 bytes, but for a WAN it can be larger or smaller.
- we must divide the datagram to make it possible for it to pass through these networks. This is called fragmentation

### Fragmented Datagram and Reassembly

- When a datagram is fragmented, each fragment has its own header.
- A fragmented datagram may itself be fragmented if it encounters a network with an even smaller MTU.
- A datagram can be fragmented by the source host or any router in the path.
- The reassembly of the datagram, however, is done only by the destination host, because each fragment becomes an independent datagram.
- The fragmented datagram can travel through different routes

### Fragmentation







# Fields Related to Fragmentation - Identification

- Identification + source IP address = uniquely define a datagram.
- When a datagram is fragmented, the value in the identification field is copied into all fragments. .
- The identification number helps the destination in reassembling the datagram.
- It knows that all fragments having the same identification value should be assembled into one datagram.

### Fields Related to Fragmentation - Flags

*The 3-bit flags field defines three flags.* 

- The leftmost bit is reserved (not used).
- The second bit (D bit) is called the do not fragment bit.
- The **third bit** (M bit) is called the more fragment bit.

### fragmentation offset

- The 13-bit fragmentation offset field shows the relative position of this fragment with respect to the whole datagram.
- It is the offset of the data in the original datagram measured in units of 8 bytes.